# QUANTUM COMPUTATION

## Lecture 3

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#### 1 Quantum entanglement (examples)

EXERCISE 1

Show that a (natural) inner product on  $H \otimes H$  is  $\langle \alpha \beta | \gamma \delta \rangle = \langle \alpha | \gamma \rangle \langle \beta | \delta \rangle$ .

EXERCISE 2

Let  $|\alpha\rangle$ ,  $|\beta\rangle$  be an orthonormal basis of  $H = \mathbb{C}^2$  and consider the observable  $X = 0 |\alpha\rangle\langle\alpha| + 1 |\beta\rangle\langle\beta|$ . Measuring  $X \otimes \mathbf{1}$  with respect to the Bell state  $|\Phi^+\rangle$  forces the second qubit to collapse to the conjugate of same state as the first qubit. Measuring with respect to  $|\Psi^{-}\rangle$  does the opposite. Try for  $|\Phi^-\rangle$  and  $|\Psi^+\rangle$ .

Solution. Let  $|\alpha\rangle = a |0\rangle + b |1\rangle$  for some  $a, b \in \mathbb{C}, |a|^2 + |b|^2 = 1$ . So  $|\beta\rangle = d |0\rangle - c |1\rangle$  such that  $(c,d) = e^{i\theta}(a^*,b^*)$  for some  $\theta$ . We have  $X = 0 |\alpha\rangle\langle\alpha| + 1 |\beta\rangle\langle\beta| = 0E_{\alpha} + 1E_{\beta}$  When we measure  $X \otimes \mathbf{1}$  in the state  $|\Phi^+\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$ , we have  $\operatorname{Prob}_{|\Phi^+\rangle}(X \otimes \mathbf{1} = 0) = \langle \Phi^+|(E_{\alpha} \otimes \mathbf{1})\Phi^+\rangle$ . Some small results:  $\langle \alpha | 0 \rangle = a^*, \langle \alpha | 1 \rangle = b^*, \langle 0 | \alpha \rangle = a, \langle 1 | \alpha \rangle = b$ . Firstly note that

$$(E_{\alpha} \otimes \mathbf{1}) |\Phi^{+}\rangle = \frac{1}{\sqrt{2}} [(a^{*} |\alpha\rangle) \otimes |0\rangle + (b^{*} |\alpha\rangle) \otimes |1\rangle]$$
  
$$= \frac{1}{\sqrt{2}} [|\alpha\rangle \otimes (a^{*} |0\rangle) + |\alpha\rangle \otimes (b^{*} |1\rangle)]$$
  
$$= \frac{1}{\sqrt{2}} |\alpha\rangle \otimes (a^{*} |0\rangle + b^{*} |1\rangle) = \frac{1}{\sqrt{2}} |\alpha\rangle \otimes |\overline{\alpha}\rangle$$

And this gives us  $\operatorname{Prob}_{|\Phi^+\rangle}(X \otimes \mathbf{1} = 0) = \langle \Phi^+ | (E_\alpha \otimes \mathbf{1}) \Phi^+ \rangle = \frac{1}{2}$ . And so  $|\tilde{\Phi}^+\rangle = |\alpha\rangle \otimes |\overline{\alpha}\rangle$ .

EXERCISE 3

Let  $|\alpha\rangle, |\beta\rangle$  be an orthonormal basis of  $H = \mathbb{C}^2$ . Then the Bell state  $|\Phi^+\rangle$  can be written as  $|\Phi^+\rangle = \frac{1}{\sqrt{2}} \left( |\alpha \overline{\alpha} \rangle + |\beta \overline{\beta} \rangle \right)$ .

Try a similar thing for the other Bell states.

#### 2 The commutator and anti-commutator

**Definition 2.1.** The *commutator* between two operators A, B is [A, B] := AB - BA.

**Definition 2.2.** The anti-commutator between two operators A, B is  $\{A, B\} := AB + BA$ .

Exercise 4

Check the following for operators A, B, C and scalars a, b:

- $1. \ [aA+bB,C] = a[A,C] + b[B,C], [C,aA+bB] = a[C,A] + b[C,B]$
- $2. \ [A,[B,C]]+[B,[C,A]]+[C,[A,B]]=0$
- 3. [A, B] = -[B, A]
- 4.  $[A, B]^* = [B^*, A^*]$
- 5. [A, BC] = [A, B]C + B[A, C], [AB, C] = A[B, C] + [A, C]B
- 6.  $AB = \frac{[A,B] + \{A,B\}}{2}$
- 7. If A, B are self-adjoint, so is i[A, B].

**Theorem 2.3** (Simultaneous diagonalization theorem). Suppose A, B are hermitian operators. Then [A, B] = 0 iff they can be simultaneously diagonalized with respect to some (common) orthonormal basis.

*Remark* 2.4. The Lie brackets used here is the quantum mechanics equivalent of the Poisson brackets used in classical mechanics. It might be a good thought exercise to figure out how the simultaneous diagonalization theorem has a physical meaning!

## 3 Heisenberg's Uncertainty Principle

**Definition 3.1.** The *expected value* or *expectation* of an observable X with respect to a state  $|\psi\rangle$  is  $\langle X \rangle_{\psi} := \langle \psi | X | \psi \rangle$ . Due to abuse of notation, we just write  $\langle X \rangle$ .

EXERCISE 5

Show that the above definition of expectation of a quantum observable is as good as the definition of expectation of a classical random variable, that is, show that if  $X = \sum_{\lambda \in \sigma(X)} \lambda E_{\lambda}$  is the spectral decomposition (all  $E_{\lambda}$ 's are rank one projections) then  $\langle \psi | X | \psi \rangle = \sum_{\lambda \in \sigma(X)} \lambda \cdot \operatorname{Prob}_{\psi}(X = \lambda)$ .

Solution.

$$\mathbb{E}(X) = \sum_{\lambda \in \sigma(X)} \lambda \cdot \operatorname{Prob}_{\psi}(X = \lambda)$$
$$= \sum_{\lambda \in \sigma(X)} \lambda \langle \psi | E_{\lambda} | \psi \rangle$$
$$= \sum_{\lambda \in \sigma(X)} \langle \psi | \lambda E_{\lambda} | \psi \rangle$$
$$= \left\langle \psi \left| \sum_{\lambda \in \sigma(X)} \lambda E_{\lambda} \right| \psi \right\rangle$$
$$= \langle \psi | X | \psi \rangle = \langle X \rangle$$

**Definition 3.2.** The standard deviation of an obvservable X with respect to the state  $|\psi\rangle$  is defined as  $\Delta_{\psi}(X) := \sqrt{\langle X^2 \rangle_{\psi} - \langle X \rangle_{\psi}^2}$ . Again due to abuse of notation, we simply write  $\Delta(X)$ .

EXERCISE 6

The above is a good definition of standard deviation, that is, verify that  $(\Delta(X))^2 = \langle (X - \langle X \rangle)^2 \rangle$ .

Solution.

$$\langle (X - \langle X \rangle)^2 \rangle = \langle X^2 - 2 \langle X \rangle X + \langle X \rangle^2 \rangle$$

$$= \left\langle \psi \Big| X^2 - 2 \langle X \rangle X + \langle X \rangle^2 \Big| \psi \right\rangle$$

$$= \left\langle \psi \Big| X^2 \Big| \psi \right\rangle - 2 \left\langle X \right\rangle \left\langle \psi | X | \psi \right\rangle + \left\langle X \right\rangle^2 \left\langle \psi | \psi \right\rangle$$

$$= \left\langle X^2 \right\rangle - 2 \left\langle X \right\rangle^2 + \left\langle X \right\rangle^2$$

$$= \left\langle X^2 \right\rangle - \left\langle X \right\rangle^2 = (\Delta(X))^2$$

**Theorem 3.3** (Heisenberg's inequality). If A, B are observables in a quantum system under the vector state  $|\rho\rangle$ , then

$$\Delta(A) \cdot \Delta(B) \ge \frac{1}{2} \left< [A, B] \right>$$

*Proof.* Let X, Y be any observables in a quantum system with state  $|\rho\rangle$ . Say  $\langle \rho | XY | \rho \rangle = a + ib$  for some  $a, b \in \mathbb{R}$ . Then we have

$$\langle \rho | [X, Y] | \rho \rangle = \langle \rho | XY | \rho \rangle - \langle \rho | YX | \rho \rangle = 2ib \langle \rho | \{X, Y\} | \rho \rangle = \langle \rho | XY | \rho \rangle + \langle \rho | YX | \rho \rangle = 2a$$

And thus,  $|\langle \rho | [X,Y] | \rho \rangle|^2 + |\langle \rho | \{X,Y\} | \rho \rangle|^2 = 4 |\langle \rho | XY | \rho \rangle|^2$  Applying the Cauchy Schwarz inequality and have,

$$\begin{split} |\langle \rho | XY | \rho \rangle|^2 &= |\langle X\rho | Y\rho \rangle|^2 \stackrel{\text{CS}}{\leq} \langle X\rho | X\rho \rangle \langle Y\rho | Y\rho \rangle = \langle \rho | X^2 | \rho \rangle \langle \rho | Y^2 | \rho \rangle = \langle X^2 \rangle \langle Y^2 \rangle \\ \implies 4 \langle X^2 \rangle \langle Y^2 \rangle \geq 4 |\langle \rho | XY | \rho \rangle|^2 = |\langle \rho | [X,Y] | \rho \rangle|^2 + |\langle \rho | \{X,Y\} | \rho \rangle|^2 \geq |\langle \rho | [X,Y] | \rho \rangle|^2 \end{split}$$

Now take  $X = A - \langle A \rangle$ ,  $Y = B - \langle B \rangle$  in the above. Check that  $\langle \rho | [X, Y] | \rho \rangle = \langle \rho | [A, B] | \rho \rangle$ . So we finally have

$$\Delta(A)\cdot \Delta(B) \geq \frac{1}{2}\left<[A,B]\right>$$

#### Quantum Gates 4

**Definition 4.1.** An *n*-qubit quantum gate is a unitary operator on  $(\mathbb{C}^2)^{\otimes n}$  (or unitary matrix, considering canonical basis).

### 4.1 1-qubit quantum gates

#### 4.1.1 Pauli gates

The NOT gate: 
$$\sigma_1 = \sigma_x = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$
.  $\sigma_1 |0\rangle = |1\rangle, \sigma_1 |1\rangle = |0\rangle$   
 $\sigma_y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}$   
 $\sigma_z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$ 

### 4.1.2 Hadamard gate

$$\begin{aligned} \mathcal{H} &:= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1\\ 1 & -1 \end{bmatrix} \\ & |+\rangle &:= \mathcal{H} |0\rangle = \frac{1}{\sqrt{2}} \left( |0\rangle + |1\rangle \right) \\ & |-\rangle &:= \mathcal{H} |1\rangle = \frac{1}{\sqrt{2}} \left( |0\rangle - |1\rangle \right) \end{aligned}$$

EXERCISE 7 Verify that  $\mathcal{H}^{\otimes n} | \boldsymbol{x} \rangle = \frac{1}{2^{n/2}} \sum_{\boldsymbol{y} \in \{0,1\}^n} (-1)^{\langle \boldsymbol{x} | \boldsymbol{y} \rangle} | \boldsymbol{y} \rangle$ 

#### 4.1.3 Phase gate

$$S := |0\rangle + i |1\rangle = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$$
  
4.1.4  $\frac{\pi}{8}$ -gate  
$$T := e^{i\frac{\pi}{8}} \begin{bmatrix} e^{-i\frac{\pi}{8}} & 0 \\ 0 & e^{i\frac{\pi}{8}} \end{bmatrix}$$

## 4.2 2-qubit quantum gates

### 4.2.1 Controlled NOT gate

This gate reads the first qubit and performs the NOT operation on the second qubit. That is, we have the following transformations:

$$\begin{array}{l} |00\rangle \mapsto |00\rangle \\ |01\rangle \mapsto |01\rangle \\ |10\rangle \mapsto |11\rangle \\ |11\rangle \mapsto |10\rangle \end{array}$$

Let  $\pi_1, \pi_2$  be the projection operators  $\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}$  (respectively) on  $\mathbb{C}^2$ . Then the *CNOT* gate can be defined as

$$CNOT := \pi_1 \otimes \mathbf{1} + \pi_2 \otimes \sigma_a$$

We can also define another CNOT gate which functions in exactly the opposite way: The control qubit is the second bit and the NOT is performed on the first bit. So define

$$C'NOT := \mathbf{1} \otimes \pi_1 + \sigma_x \otimes \pi_2$$

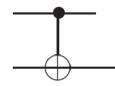


Figure 1: CNOT gate (with control in the top channel)

The matrix forms look like

$$CNOT = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \qquad \qquad C'NOT = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

EXAMPLE. The column matrix forms of the vectors  $|00\rangle, |01\rangle, |10\rangle, |11\rangle \in H \otimes H$  would be  $\begin{bmatrix} 0 \end{bmatrix}$  $\left|, \begin{array}{c} 0\\ 0\\ 0\\ 1 \end{array}\right| \text{ respectively. So, applying } CNOT \text{ on the vector } \left(a \left|0\right\rangle + b \left|1\right\rangle\right) \otimes \left|0\right\rangle \text{ gives }\right|$  $\begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix}$ 1

$$(CNOT)(a|00\rangle + b|10\rangle) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} a \\ 0 \\ b \\ 0 \end{bmatrix} = \begin{bmatrix} a \\ 0 \\ b \\ b \end{bmatrix} = a|00\rangle + b|11\rangle$$

#### 4.2.2 Controlled U gate

This gate is similar to the controlled NOT gate. The first qubit acts as the control bit, and the unitary operation U is performed on the second qubit. That is, we have the following transformations:

$$\begin{array}{l} |00\rangle \mapsto |0\rangle \otimes |0\rangle \\ |01\rangle \mapsto |0\rangle \otimes |1\rangle \\ |10\rangle \mapsto |1\rangle \otimes (U \, |0\rangle) \\ |11\rangle \mapsto |1\rangle \otimes (U \, |1\rangle) \end{array}$$
The gate can be defined by
$$CU := \pi_1 \otimes \mathbf{1} + \pi_2 \otimes U$$
Similarly we can define
$$C'U := \mathbf{1} \otimes \pi_1 + U \otimes \pi_2$$

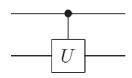


Figure 2: Controlled U gate (with control in the top channel)

The matrix form looks like

$$CU = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & a & b \\ 0 & 0 & c & d \end{bmatrix} \qquad \qquad C'U = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & a & 0 & b \\ 0 & 0 & 1 & 0 \\ 0 & c & 0 & d \end{bmatrix}$$

where  $U = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$ .

Remark 4.2. The CNOT and C'NOT are just the cases when  $U = \sigma_x$ .

### 4.2.3 SWAP gate

The SWAP gate does as you might expect: it swaps the two qubits. So the transformations are as follows:

$$\begin{array}{l} |00\rangle \mapsto |00\rangle \\ |01\rangle \mapsto |10\rangle \\ |10\rangle \mapsto |01\rangle \\ |11\rangle \mapsto |11\rangle \end{array}$$

It is relatively easier to construct the matrix first

$$SWAP = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

EXERCISE 8

Verify that SWAP = (CNOT)(C'NOT)(CNOT)

The above exercise helps us to define

$$SWAP := \pi_1 \otimes \pi_1 + \pi_2 \otimes \pi_1 + (X\pi_1) \otimes (X\pi_2) + (X\pi_2) \otimes (X\pi_1)$$

### 4.3 3-qubit quantum gates

## 4.3.1 Toffoli Gate (CCNOT)

 $CCNOT = \begin{bmatrix} \mathbf{1}_6 & \mathbf{0}_{6\times 2} \\ \mathbf{0}_{2\times 6} & X \end{bmatrix}$ 



Figure 3: The Toffoli Gate (or *CCNOT* gate) with control in the first two qubits

## 4.4 Preparing Bell states

Bell states can be prepared by the action of  $\mathcal{H} \otimes \mathbf{1}$  followed by CNOT on the standard bases of  $\mathbb{C}^2 \otimes \mathbb{C}^2$ .

$$(CNOT)(\mathcal{H} \otimes \mathbf{1}) |00\rangle = CNOT \left(\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \otimes |0\rangle\right)$$
$$= CNOT \left(\frac{1}{\sqrt{2}}(|00\rangle + |01\rangle)\right)$$
$$= \frac{1}{\sqrt{2}}(|00\rangle + |01\rangle)$$