

# Fiedler Vector Method

Project for Linear algebra and its Applications

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- 1 The problem
- 2 Some spectral Graph Theory
- 3 What exactly do we want?
- 4 Where does spectral graph theory come in?
- 5 Our new (approximate) problem and solution
- 6 Why does this give connected subgraphs?
- 7 Some examples
- 8 Some questions

## Statement

Given a (finite simple connected) graph  $G = (V, E)$ , partition  $V = V_1 \sqcup V_2$  such that the number of “cut” edges is minimized, while keeping  $|V_1| \simeq |V_2|$ .

- 1 The problem
- 2 Some spectral Graph Theory
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- 8 Some questions

# Terminology and definitions

To every finite undirected graph  $G = (V, E)$ , we can associate the following matrices:

- *Adjacency matrix*  $A$  of size  $|V| \times |V|$ .  $A_{i,j} = \mathbf{1}_{\{i,j\} \in E}$ .
- *Degree matrix*  $D$  of size  $|V| \times |V|$ .  $D_{i,j} = \delta_i^j \cdot \deg(i)$ .
- *Incidence matrix*  $H$  of size  $|V| \times |E|$ .  $H_{v,e} = \begin{cases} 1 & \text{if } v \text{ adjacent to edge } e \\ 0 & \text{else} \end{cases}$ .
- *Laplacian*  $\mathcal{L}$  of size  $|V| \times |V|$ .  $L = D - A$ .

# Elementary results

$\mathcal{L}$  has real eigenvalues.

$\mathcal{L}$  is real symmetric. ■

**Fact:**  $\mathcal{L} = HH^t$ .

$\mathcal{L}$  is positive semidefinite.

$\mathcal{L}\mathbf{v} = \lambda\mathbf{v}$ , then  $\lambda\|\mathbf{v}\|^2 = \langle \mathcal{L}\mathbf{v}, \mathbf{v} \rangle = \mathbf{v}^t HH^t \mathbf{v} = \|H^t \mathbf{v}\|^2$ . ■

$\mathcal{L}$  has only nonnegative eigenvalues. We note that the eigenvalue 0 is always achieved. Indeed,  $\mathcal{L}(1, \dots, 1) = (0, \dots, 0)$ .

- 1 The problem
- 2 Some spectral Graph Theory
- 3 What exactly do we want?**
- 4 Where does spectral graph theory come in?
- 5 Our new (approximate) problem and solution
- 6 Why does this give connected subgraphs?
- 7 Some examples
- 8 Some questions

# Modelling a partition

Recall that a partition of graph  $G = (V, E)$  is a partition of  $V = V_1 \sqcup V_2$ . For a partition  $P = (V_1, V_2)$ , define its cut  $\mathfrak{C}(P)$  to be the number of edges in  $E$  with at least one endpoint in each  $V_i$ .

Think of a partition as assigning  $\pm 1$  to each vertex, with the agreement that vertex  $i$  lies in  $V_1$  when it's assigned  $+1$ , and it's in  $V_2$  otherwise. Say vertex  $i$  is assigned  $x_i$ , which gives a vector  $\mathbf{x} = (x_1, \dots, x_n)$ .



# What is the cut value for the aforesaid assignment?

The cut value of partition  $P = (V_1, V_2)$  is  $\mathfrak{C}(P) = \sum_{\substack{\{i,j\} \in E \\ x_i \neq x_j}} 1$ . Some clever manipulation gives

$$\begin{aligned}
 4 \cdot \mathfrak{C}(P) &= 4 \sum_{\substack{\{i,j\} \in E \\ x_i \neq x_j}} 1 = \sum_{\substack{\{i,j\} \in E \\ x_i \neq x_j}} (\pm 2)^2 + \sum_{\substack{\{i,j\} \in E \\ x_i = x_j}} 0^2 \\
 &= \sum_{\substack{\{i,j\} \in E \\ x_i \neq x_j}} (x_i - x_j)^2 + \sum_{\substack{\{i,j\} \in E \\ x_i = x_j}} (x_i - x_j)^2 \\
 &= \sum_{\{i,j\} \in E} (x_i - x_j)^2
 \end{aligned}$$

# What does $|V_1| \simeq |V_2|$ mean?

If  $V_1$  and  $V_2$  were to have the same number of vertices, it simply means that equal number of vertices are assigned  $+1$  and  $-1$ . This is to say that  $\sum x_i = 0$ .

Conversely  $\sum x_i = 0$  (where each  $x_i \in \{\pm 1\}$ ) means that same number of vertices are assigned  $\pm 1$  each. That simply means  $|V_1| = |V_2|$ .

Now  $|V_1| \simeq |V_2|$  simply means we want  $|\sum x_i|$  to be as small as possible. That is,  $\sum x_i \simeq 0$ .

## GOAL

Find  $\mathbf{x} \in \{\pm 1\}^n$  with  $\sum x_i \simeq 0$  such that

$$\sum_{\{i,j\} \in E} (x_i - x_j)^2$$

is minimized.

Note that the condition of each  $x_i$  being  $\pm 1$  implies  $\sum x_i^2 = n$ . The converse is true if we ask each  $x_i \in [-1, 1]$ . So we are not giving up much.

Asking for only an approximate solution, we replace the condition  $\sum x_i \simeq 0$  with  $\sum x_i = 0$ .

# NEW (approximate) GOAL

Find  $\mathbf{x} \in \mathbb{R}^n$  with  $\sum x_i = 0$  such that

$$\sum_{\{i,j\} \in E} (x_i - x_j)^2$$

is minimized (subject to  $\sum_i x_i^2 = 1$ ).

- 1 The problem
- 2 Some spectral Graph Theory
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- 8 Some questions

## Bring in the matrices ...

Some clever manipulation:

$$\begin{aligned} \sum_{\{i,j\} \in E} (x_i - x_j)^2 &= \sum_{\{i,j\} \in E} (x_i^2 + x_j^2 - 2x_i x_j) = \sum_i \underbrace{\deg(i)}_{d_i} x_i^2 - 2 \sum_{\{i,j\} \in E} x_i x_j \\ &= \mathbf{x}^t \begin{bmatrix} d_1 & & \\ & \ddots & \\ & & d_n \end{bmatrix} \mathbf{x} - \mathbf{x}^t A \mathbf{x} = \mathbf{x}^t (D - A) \mathbf{x} = \mathbf{x}^t \mathcal{L} \mathbf{x} \end{aligned}$$

Illustration for =

$$\begin{bmatrix} x & y \end{bmatrix} \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} x & y \end{bmatrix} \begin{bmatrix} a_{11}x + a_{12}y \\ a_{21}x + a_{22}y \end{bmatrix} = a_{11}x^2 + (a_{12} + a_{21})xy + a_{22}y^2.$$

If the above matrix is symmetric, then the result is  $a_{11}x^2 + \underline{2a_{12}}xy + a_{22}y^2$ .

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- 2 Some spectral Graph Theory
- 3 What exactly do we want?
- 4 Where does spectral graph theory come in?
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- 7 Some examples
- 8 Some questions

So our problem becomes . . .

For what  $\mathbf{x}$  is  $\mathbf{x}^t \mathcal{L} \mathbf{x}$  minimized, subject to  $\|\mathbf{x}\|^2 = n$  and  $\langle \mathbf{x}, (1, \dots, 1) \rangle = 0$ .

An equivalent optimization problem is

Find the argument  $\mathbf{x}$  for

$$\min_{\substack{\mathbf{x} \perp (1, \dots, 1) \\ \|\mathbf{x}\| \neq 0}} \frac{\mathbf{x}^t \mathcal{L} \mathbf{x}}{\|\mathbf{x}\|^2} = \min_{\substack{\mathbf{x} \perp (1, \dots, 1) \\ \|\mathbf{x}\| \neq 0}} R_{\mathcal{L}}(\mathbf{x}).$$



# Recall from class

If  $A_{n \times n}$  is a real symmetric positive semidefinite matrix with eigenvalues  $0 \leq \lambda_1 \leq \dots \leq \lambda_n$  with eigenbasis  $(\mathbf{v}_1, \dots, \mathbf{v}_n)$  (such that  $A\mathbf{v}_i = \lambda_i\mathbf{v}_i$ ) then

$$\min_{\substack{\mathbf{x} \perp \langle \mathbf{v}_1, \dots, \mathbf{v}_k \rangle \\ \|\mathbf{x}\| \neq 0}} R_{\mathcal{L}}(\mathbf{x}) = \lambda_{k+1}.$$

# Our interest is

In particular for  $k = 1$

$$\min_{\substack{\mathbf{x} \perp (1, \dots, 1) \\ \|\mathbf{x}\| \neq 0}} R_{\mathcal{L}}(\mathbf{x}) = \lambda_2$$

because 0 is always an eigenvalue of  $\mathcal{L}$  corresponding to the eigenvector  $(1, 1, \dots, 1)$ .

# Our algorithm

Initially we had asked that we'll take  $i \in V_1$  if  $x_i = +1$ , and  $i \in V_2$  if  $x_i = -1$ . But while arriving to our optimization problem, we were 'loose' about  $x_i \in \{\pm 1\}$ , we now our algorithm becomes:

Look at the  $i^{\text{th}}$  component of a Fiedler vector. If it is  $< 0$ , put the  $i^{\text{th}}$  vertex in  $V_1$ , otherwise put it in  $V_2$ .

- 1 The problem
- 2 Some spectral Graph Theory
- 3 What exactly do we want?
- 4 Where does spectral graph theory come in?
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- 8 Some questions

### Theorem (Fiedler's theorem of connectivity of spectral graph partitions)

*Let  $G = (V, E)$  be a connected graph and let  $\mathbf{x} = (x_1, \dots, x_n)$  is the Fiedler vector for the Laplacian of this graph. Let  $V_1 = \{i : x_i > 0\}$  and  $V_2 = V \setminus V_1$ . Let  $G_i$  be the subgraphs induced by  $V_i$ . Then  $G_i$  are both connected.*

## Suppose not ...

(Re)label the vertices in a way such that  $V_1 = \{1, \dots, k\}$ ,  $V_2 = \{k + 1, \dots, n\}$ . For the sake of contradiction suppose  $G_1$  has (at least) two connected components; say the vertices are  $\{1, \dots, t\}$  and  $\{t + 1, \dots, k\}$ . So the Laplacian looks like

$$\mathcal{L} = \begin{bmatrix} L_{11} & \mathbf{0} & L_{13} \\ \mathbf{0} & L_{22} & L_{23} \\ L_{13}^T & L_{23}^T & L_{33} \end{bmatrix}.$$

By nature of  $\mathcal{L}$ , the entries of  $L_{13}, L_{23}$  are nonpositive. Write the similar block form for

$$\mathbf{x} = \begin{bmatrix} (\mathbf{x}_1)_{t \times 1} \\ (\mathbf{x}_2)_{(k-t) \times 1} \\ (\mathbf{x}_3)_{(n-k) \times 1} \end{bmatrix}.$$

where components of  $\mathbf{x}_1, \mathbf{x}_2$  are positive and those of  $\mathbf{x}_3$  are negative.

We produce two eigenvalues of  $\mathcal{L}$  which are less than  $\lambda_2$ . Contradiction!

$$\mathcal{L}\mathbf{x} = \lambda_2\mathbf{x} \implies L_{11}\mathbf{x}_1 + L_{13}\mathbf{x}_3 = \lambda_2\mathbf{x}_1$$

Fix any positive real  $\varepsilon > 0$ . Then  $(\varepsilon I + L_{11})\mathbf{x}_1 + L_{13}\mathbf{x}_3 = (\varepsilon + \lambda_2)\mathbf{x}_1$ .

Assume that  $(\varepsilon I + L_{11})$  is invertible and its inverse  $Y$  is a positive matrix. Multiplying both sides of the above equation by  $Y$ , we get

$$\begin{aligned} \mathbf{x}_1 + YL_{13}\mathbf{x}_3 &= (\varepsilon + \lambda_2)Y\mathbf{x}_1 \\ \implies \mathbf{x}_1^T\mathbf{x}_1 + \mathbf{x}_1^TYL_{13}\mathbf{x}_3 &= (\varepsilon + \lambda_2)\mathbf{x}_1^TY\mathbf{x}_1 \\ \implies (\varepsilon + \lambda_2)\frac{\mathbf{x}_1^TY\mathbf{x}_1}{\mathbf{x}_1^T\mathbf{x}_1} &= 1 + \frac{\mathbf{x}_1^TYL_{13}\mathbf{x}_3}{\mathbf{x}_1^T\mathbf{x}_1} > 1 \\ \implies (\varepsilon + \lambda_2)\lambda_t(Y) &= (\varepsilon + \lambda_2)\max_{v \neq 0} \frac{v^TYv}{v^Tv} > 1 \\ \implies \lambda_1(\varepsilon I + L_{11}) &= 1/\lambda_t(Y) < \varepsilon + \lambda_2 \\ \implies \varepsilon + \lambda_1(L_{11}) &< \varepsilon + \lambda_2 \end{aligned}$$

Hence  $\lambda_1(L_{11}) < \lambda_2$ . Similarly,  $\lambda_1(L_{22}) < \lambda_2$ .

If the eigenvectors corresponding to  $\lambda_1(L_{11})$  and  $\lambda_1(L_{22})$  are  $v_1$  and  $v_2$  respectively, then

$$\begin{bmatrix} L_{11} & 0 \\ 0 & L_{22} \end{bmatrix} \begin{bmatrix} v_1 \\ 0 \end{bmatrix} = \lambda_1(L_{11}) \begin{bmatrix} v_1 \\ 0 \end{bmatrix} \quad \text{and} \quad \begin{bmatrix} L_{11} & 0 \\ 0 & L_{22} \end{bmatrix} \begin{bmatrix} v_2 \\ 0 \end{bmatrix} = \lambda_1(L_{22}) \begin{bmatrix} v_2 \\ 0 \end{bmatrix} .$$

Thus the matrix  $\begin{bmatrix} L_{11} & 0 \\ 0 & L_{22} \end{bmatrix}$  has two eigenvalues less than  $\lambda_2$ .

### Theorem (Cauchy Interlacing theorem)

*Let  $A_{n \times n}$  be a symmetric matrix and  $B_{m \times m}$  be a principal submatrix of  $A$ . Further, let the eigenvalues of  $A$  be  $\lambda_1 \leq \dots \leq \lambda_n$  and the eigenvalues of  $B$  be  $\beta_1 \leq \dots \leq \beta_m$ . Then, for all  $k \leq m$ , the matrix  $A$  has at least  $k$  eigenvalues less than or equal to  $\beta_k$ .*

By Cauchy Interlacing theorem,  $\mathcal{L}$  has two eigenvalues less than  $\lambda_2$ , which is a contradiction!



Now we show that  $(\varepsilon I + L_{11})$  is invertible and its inverse  $Y$  is positive.

$$\begin{aligned} & \varepsilon I + L_{11} \\ &= D - N \\ &= D^{1/2}(I - D^{-1/2}ND^{-1/2})D^{1/2} \\ &= D^{1/2}(I - M)D^{1/2} \end{aligned}$$

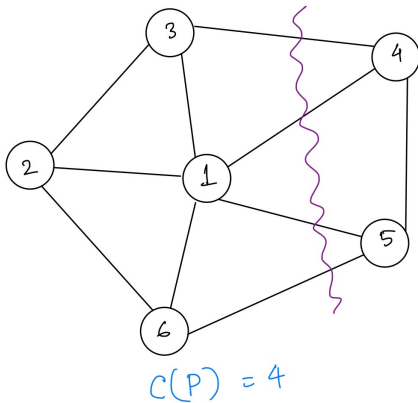
Because of some useful properties of  $M$ , it can be shown that  $(I - M)^{-1} = \sum_{l=0}^{\infty} M^l$ .

$$\begin{aligned} \therefore Y &= (\varepsilon I + L_{11})^{-1} \\ &= D^{-1/2}(I - M)^{-1}D^{-1/2} \\ &= D^{-1/2} \cdot \sum_{l=0}^{\infty} M^l \cdot D^{-1/2} \end{aligned}$$

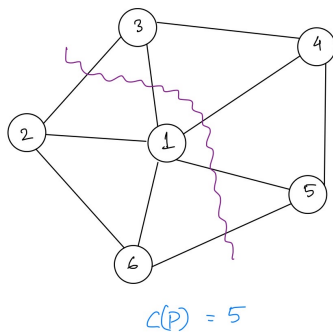
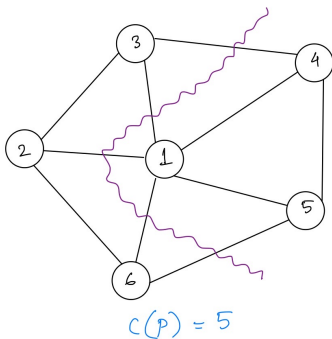
$Y$  is positive as  $\sum_{l=0}^{\infty} M^l$  is positive.

- 1 The problem
- 2 Some spectral Graph Theory
- 3 What exactly do we want?
- 4 Where does spectral graph theory come in?
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- 6 Why does this give connected subgraphs?
- 7 Some examples**
- 8 Some questions

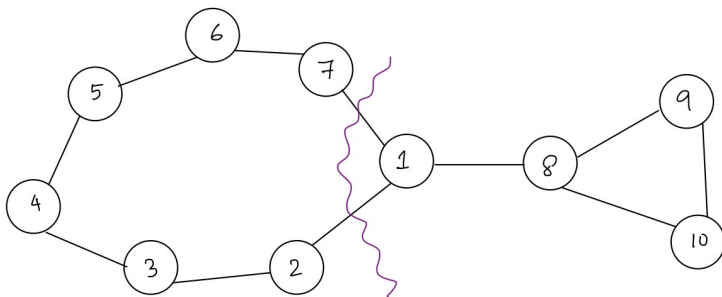
# Example 1



## Example 1: But...

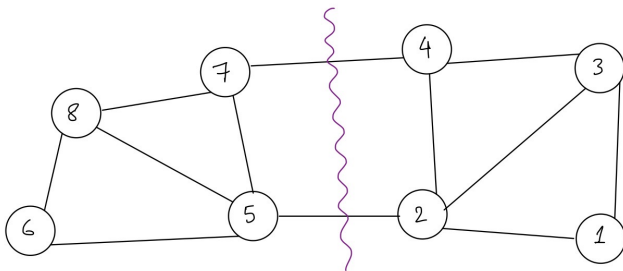


## Example 2



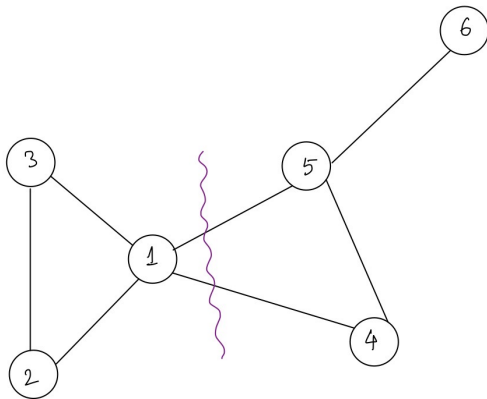
$$C(P) = 2$$

## Example 3



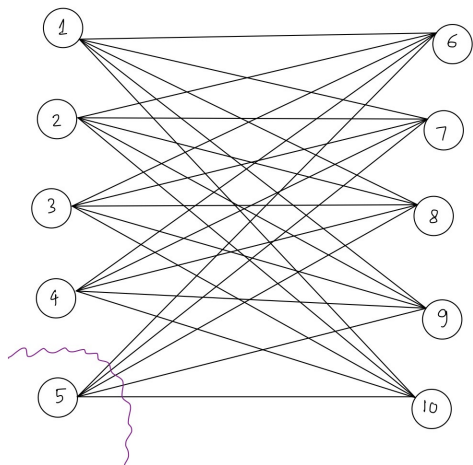
$$c(P) = 2$$

## Example 4



$$c(P) = 2$$

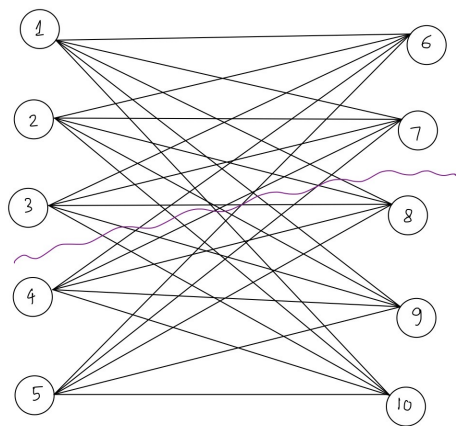
## Example 5



$$C(P) = 5$$



## Example 5: But...



$$C(P) = 13$$

- 1 The problem
- 2 Some spectral Graph Theory
- 3 What exactly do we want?
- 4 Where does spectral graph theory come in?
- 5 Our new (approximate) problem and solution
- 6 Why does this give connected subgraphs?
- 7 Some examples
- 8 Some questions**

# What effect does connectedness have on $\lambda_2$ ?

## Lemma

$\lambda_2 > 0 \iff$  the graph connected.

## Proof.

If  $G$  has two connected components, then  $(1 \cdots, 1, 0 \cdots, 0)$  and  $(1 \cdots, 1, 0 \cdots, 0)$  are LI eigenvectors for 0.

If  $G$  is connected, then use  $\mathbf{v}^t \mathcal{L} \mathbf{v} = \sum_{\{i,j\} \in E} (v_i - v_j)^2$  to show that geometric (=algebraic) multiplicity of 0 is one. ■

## Terminology

$\lambda_2$  is also called the algebraic connectivity of the graph.

# What if all components are positive?

Say  $\mathcal{L}\mathbf{v} = \lambda\mathbf{v}$  with  $\lambda > 0$ .

$$l_{11}v_1 + \cdots + l_{1n}v_n = \lambda v_1$$

$$\vdots$$

$$l_{n1}v_1 + \cdots + l_{nn}v_n = \lambda v_n$$

Adding these gives  $\sum_{i=1}^n l_{i1}v_1 + \cdots + \sum_{i=1}^n l_{in}v_n = \lambda \sum_{i=1}^n v_i$ . All the red sums are 0 because of the way  $\mathcal{L}$  is defined. It follows that  $\sum_i v_i = 0$  because  $\lambda > 0$ .

## How to balance so many 0's (if at all)?

What do we do if there are only a few nonzero components of the Fiedler vector and a bunch of 0's?

One conjecture we made was

For a connected graph  $G$  with Laplacian  $\mathcal{L}$ , let  $\mathbf{v}$  be a Fiedler vector. Look at  $S = \{i : v_i = 0\}$ . Then for each  $i \in S$ ,  $\exists j, k$  such that  $v_j > 0, v_k < 0$  and both  $j$  and  $k$  are neighbours of  $i$ .

The above conjecture is false. A counterexample is  $K_{n,n}$ , the complete bipartite graph on  $2n$  vertices.